# TreeLine: An Update-In-Place Key-Value Store for Modern Storage

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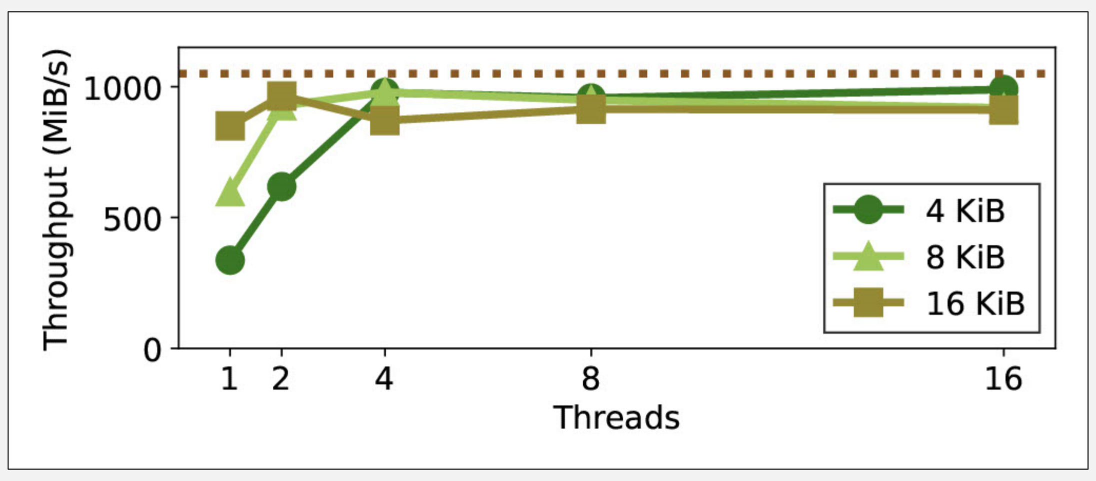
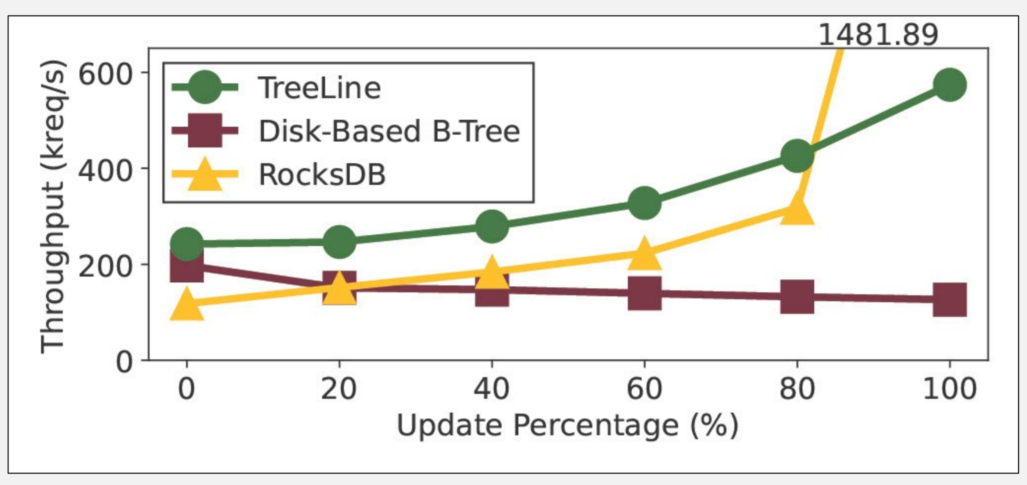
## Motivation




Both experiments on an Intel DC P4510 NVMe SSD.

- Modern persistent key-value stores, such as RocksDB [1] and LevelDB [2], typically use log-structured merge trees (LSMs) [3].
- Stellar write performance: large sequential writes exploit the high sequential throughput.
- Slow read performance: need caches, Bloom filters [4,5], compaction strategies [6,7,8]—complex and hard-to-tune. [9]

### Random Writes ≈ Sequential Writes on NVMe SSDs

- Random write throughput across (i) request sizes, and (ii) number of writing threads.
- With high parallelism, can reach advertised peak sequential write throughput [10].

### LSMs Leave Read Performance on the Table

- Zipfian-distributed ($\theta$ = 0.79) workload of reads, updates, and scans on 64 B records.
- Of the requests that are not updates, 10% are range scans and the rest are point reads.
- For read-heavy, the disk-based B-tree outperforms/is competitive against RocksDB.
- TreeLine can outperform both systems all the way up to 80% updates.

## Design Overview

### Key Idea A: Record Caching

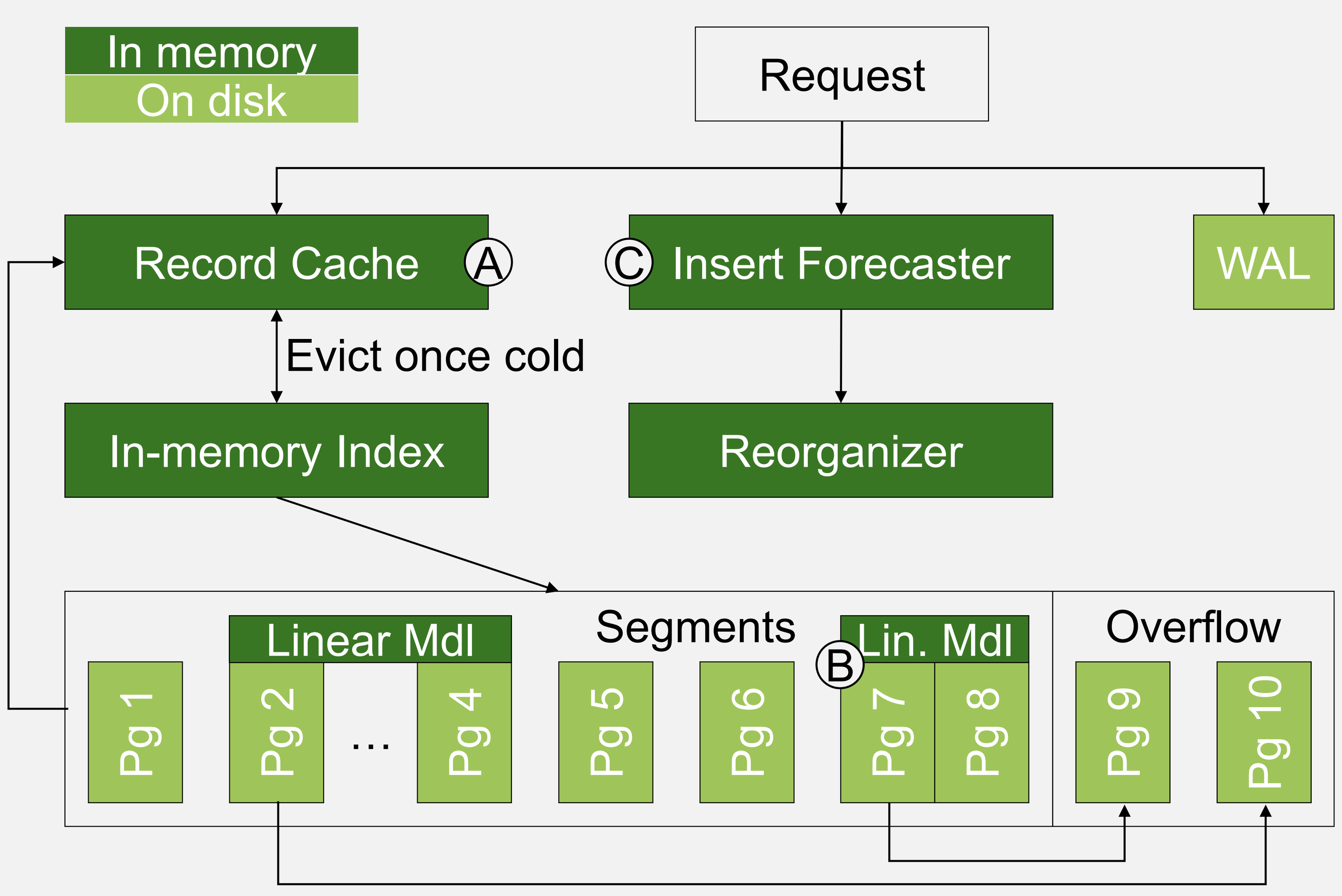
- Key-to-page mapping is expensive to change in update-in-place design.
- But variable hotness among records on the same page.
- **Solution:** Only cache records, to increase memory efficiency.

### Key Idea B: Page Grouping

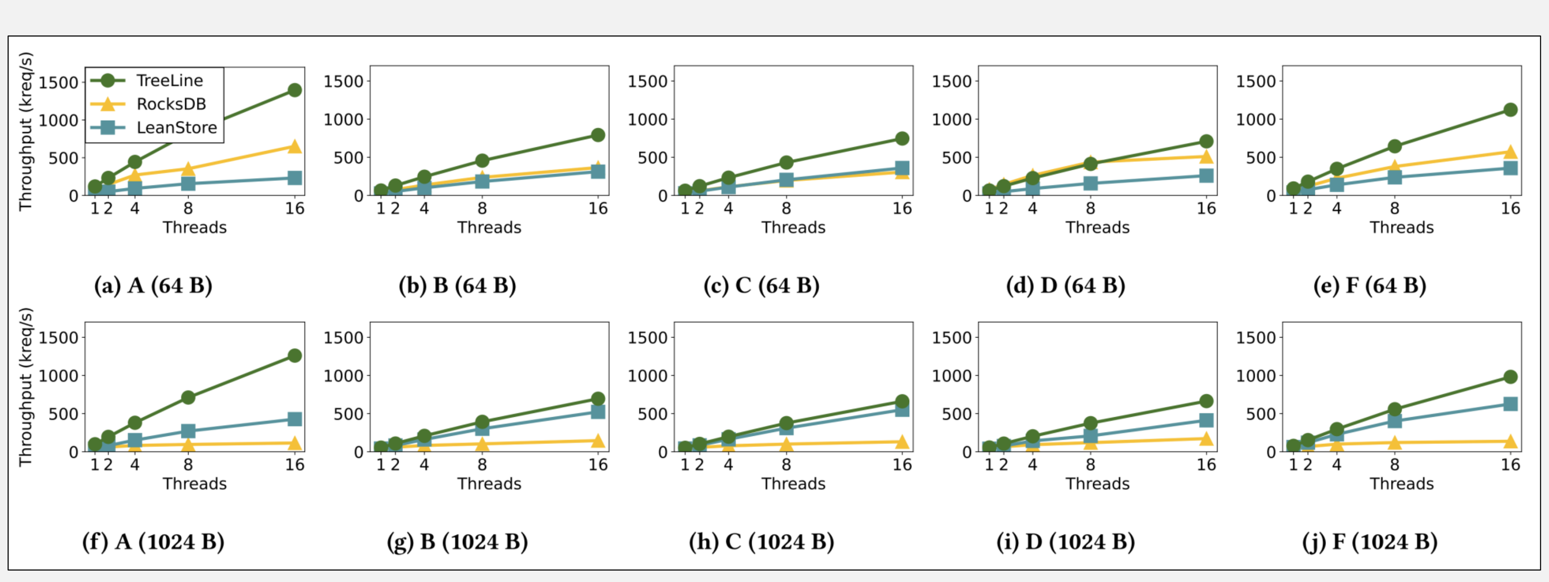
- Unlike writes, sequential reads still faster vs. random reads on modern SSDs.
- Pages must be sequential to make scans fast.
- **Solution:** Group pages into contiguous segments. Use linear models to index records within segments.

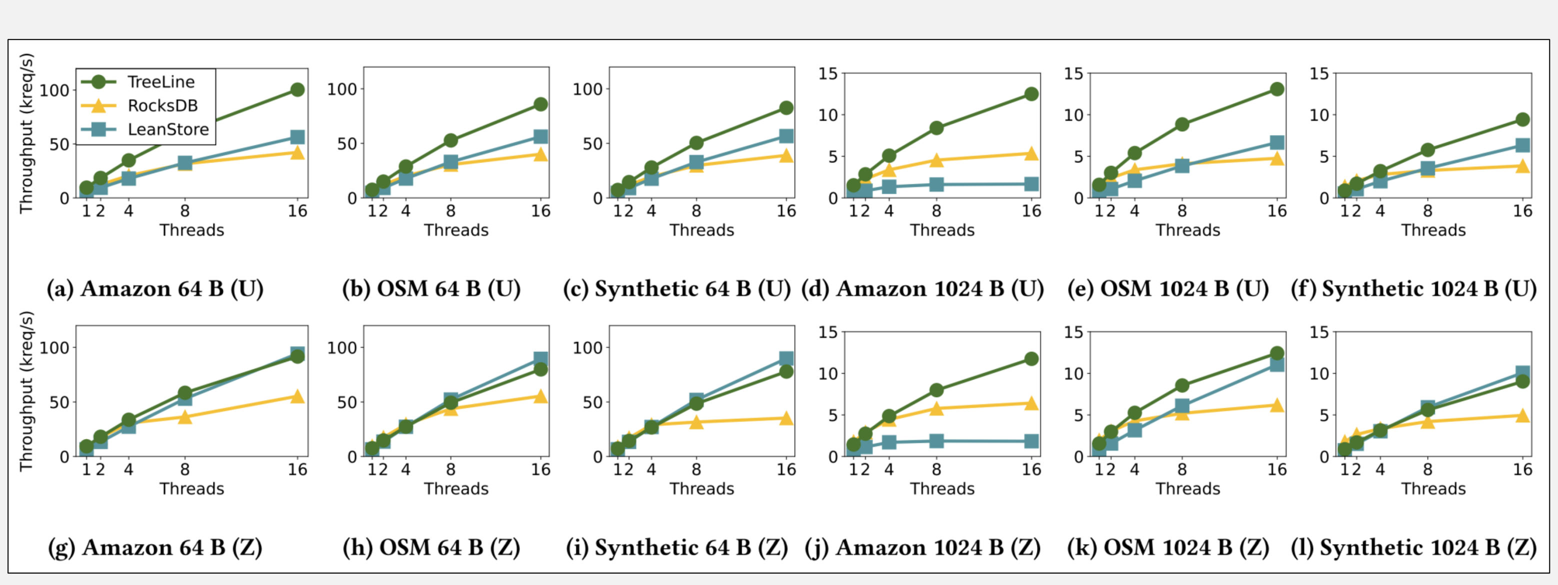
### Key Idea C: Insert Forecasting

- To avoid constant reorganization, pages should have some empty space.
- But too much empty space increases I/O amplification.
- **Solution:** Leave empty space based on epoch-based insert forecast.




## Evaluation



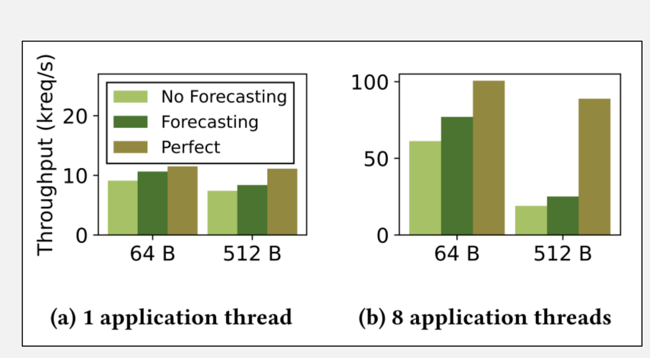




### Point Workloads

- TreeLine outperforms RocksDB (LeanStore) by **1.62× (2.81×)** and **2.99× (1.53×)** on average for **64 B**/**1024 B** records.

### Scans

- With 16 request threads, for uniform scans, TreeLine outperforms RocksDB (LeanStore) by **2.21× (1.58×)** and **2.50× (2.80×)** on average for **64 B**/**1024 B** records.
- With 16 request threads, for Zipfian-distributed scans, TreeLine outperforms RocksDB (LeanStore) by **1.74× (0.91×)** and **1.88× (1.86×)** on average for **64 B**/**1024 B** records.

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